cREXX Progress Update

The 32nd Annual Rexx Symposium

Adrian Sutherland • 8.11.2021 (Final)

cREXX Progress Update

Reminder - what I said last year

cREXX Architecture and Status

cREXX in Action

How to Help?

Thanks!

Next Up - Rexx in the RexxLA Website with Mark Hessling

Typical Bytecode Optimisation

- Threaded code
- 2. Super-instructions / inlining
- 3. Top-of-stack in a register
- 4. Scheduling the dispatch of the next VM instruction

In all about 2x faster than classic bytecode

We should be aiming for performance of only 2-5 times slower than native code

```
char code[] = {
                                                 void *code[] = {
 ICONST 1, ICONST 2,
                                                   &&ICONST 1, &&ICONST 2,
  IADD, ...
                                                   &&IADD, ...
char *pc = code;
                                                 void **pc = code;
/* dispatch Loop */
                                                 /* implementations */
while(true) {
                                                 goto **(pc);
  switch(*pc++) {
  case ICONST 1: *++sp = 1; break;
                                                 ICONST 1: pc++; *++sp = 1; goto **(pc);
  case ICONST 2: *++sp = 2; break;
                                                 ICONST 2: pc++; *++sp = 2; goto **(pc);
  case IADD:
                                                 IADD:
    sp[-1] += *sp; --sp; break;
                                                   pc++; sp[-1] += *sp; --sp; goto **(pc);
  . . .
}}
```

Pure Bytecode

Threaded Interpreter

NOTE - We could be talking about any language ...

REXX Assembler

This is where optimisations become REXX specific ...

BREXX

- Stack Based
- Leaves work to the interpreter

CREXX

- Register Based
- Trying to handle REXXisms at the low level

NEWCLAUSE			in: locals=3 {r1="A", r2="B"}
CREATE	"A"	ILOAD	r1,10
PUSH	10		
COPY			
NEWCLAUSE		ILOAD	r2,5
CREATE	"B"		
PUSH	5		
COPY			
NEWCLAUSE		IADD	r3,r1,r2
PUSHTMP			
LOAD	"A"		
LOAD	"B"		
ADD			
SAY			
NEWCLAUSE		ISAY	r3
IEXIT		HALT	
BREXX		CREXX	

We need to get this right for LLVM ...

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REXX Variable Types

- Rexx is typeless ... and more than that conceptually all variables are strings
- 2. Rexx stems provide a flexible and arbitrary index scheme
- VALUE(), INTERPRET(), and REXXSAA/EXECOMM all require dynamic variable name resolution
- Performance requires compile time resolution of variable names and types, wherever possible

Mapped to a Register when needed

Dynamic lookup (tree/ hash)

CREXX Variable

Variable Status (what buffers are uptodate)

String Buffer

Integer Buffer

Arbitrary Precision Number
Buffer

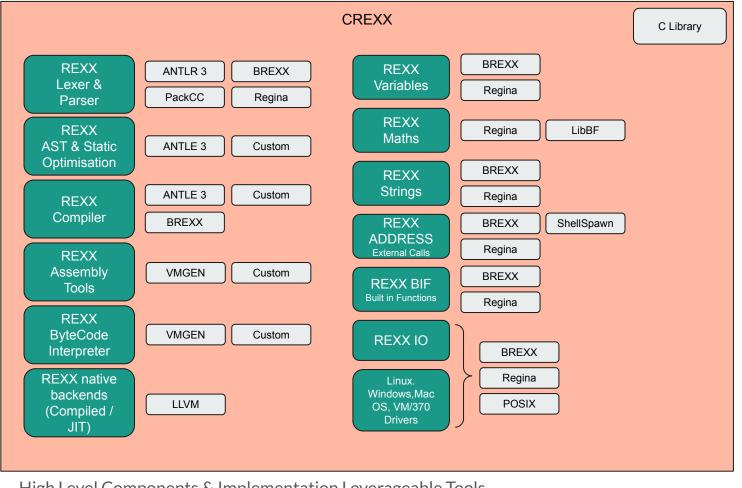
Date Buffer ...

Note: BREXX has a single shared buffer

Lazy updates (only updates the type being set)

Copy on Write

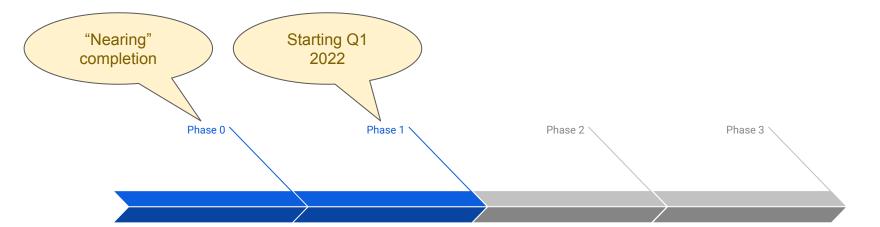
REXX VM Specific Memory Allocation



High Level Components & Implementation Leverageable Tools



cREXX Architecture and Status



Proof of Concept

Goal: Sustainability

Prove architectural concepts and the ability for the project to deliver by creating a modern REXX implementation

Classic REXX

Goal: Standards compliancy

Formalise the implementation by creating a high quality, stable, performantand compliant Classic REXX

Native Performance

Goal: Native Binaries

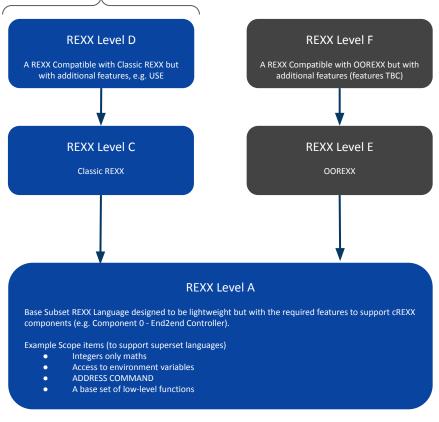
Integrate to the LLVM backend to allow optimised native binaries for multiple target operating systems

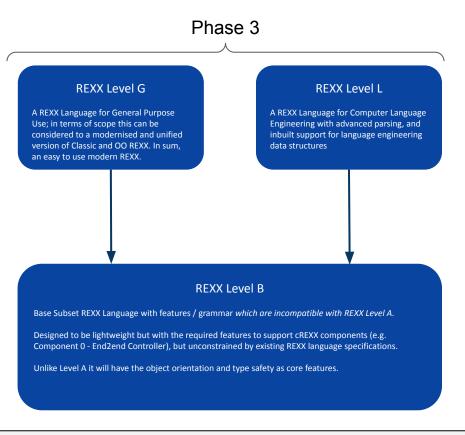
REXX Modernisation

Goal: Contemporary REXX

Re-imagine REXX for new users and workloads, and with contemporary language features

Phases 0 to 2





Superset REXX Level In Project Scope Scope TBC

Phases 0 to 2 Phase 3 **REXX Level D REXX Level F** A REXX Compatible with OOREXX but with A REXX Compatible with Classic REXX but with additional features, e.g. USE additional features (features TBC) **REXX Level G REXX Level L** A REXX Language for General Purpose A REXX Language for Computer Language Use; in terms of scope this can be Engineering with advanced parsing, and considered to a modernised and unified inbuilt support for language engineering version of Classic and OO REXX. In sum, **REXX Level C REXX Level E** data structures an easy to use modern REXX. Classic REXX OOREXX REXX Level B (Phase 0) Base Subset REXX Language with features / grammar which are incompatible with REXX Level C. Designed to be lightweight but with the required features to support cREXX components (e.g. Component 0 - End2end Controller), but unconstrained by existing REXX language specifications. Access to environment variables ADDRESS COMMAND A base set of low-level functions (via ASSEMBLE instruction) It will have the object orientation and type safety as core features.

Superset

REXX Level

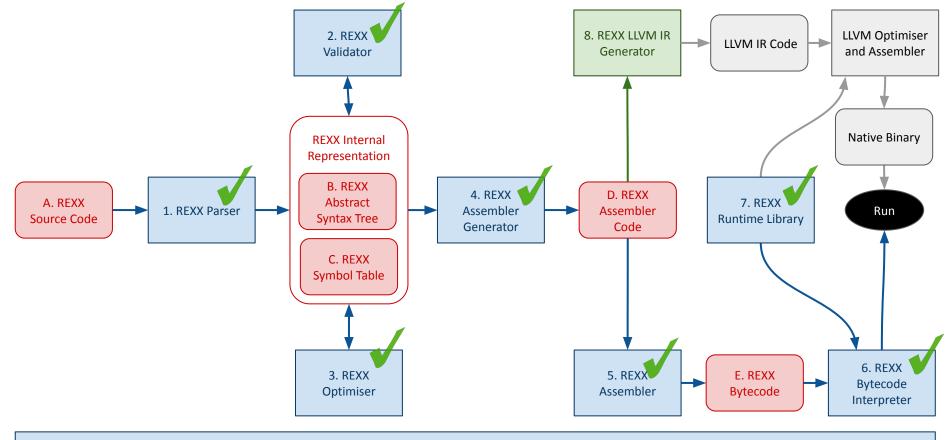
Key

Subset REXX

Level

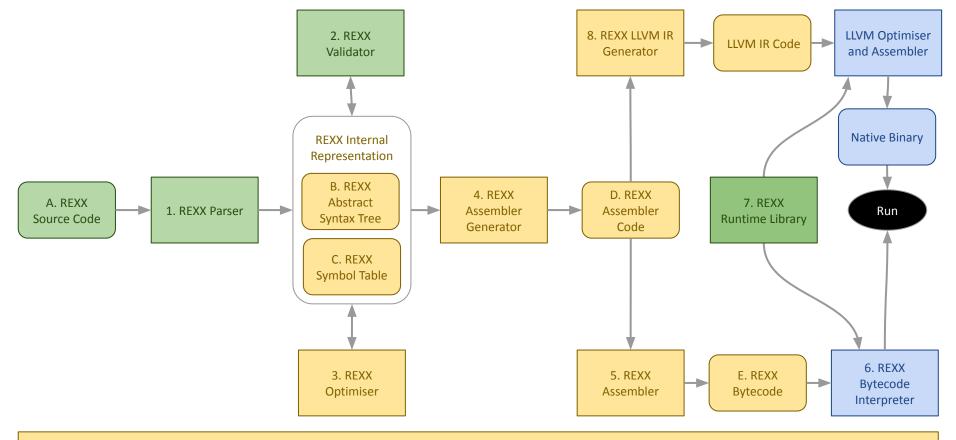
In Project Scope

Scope TBC

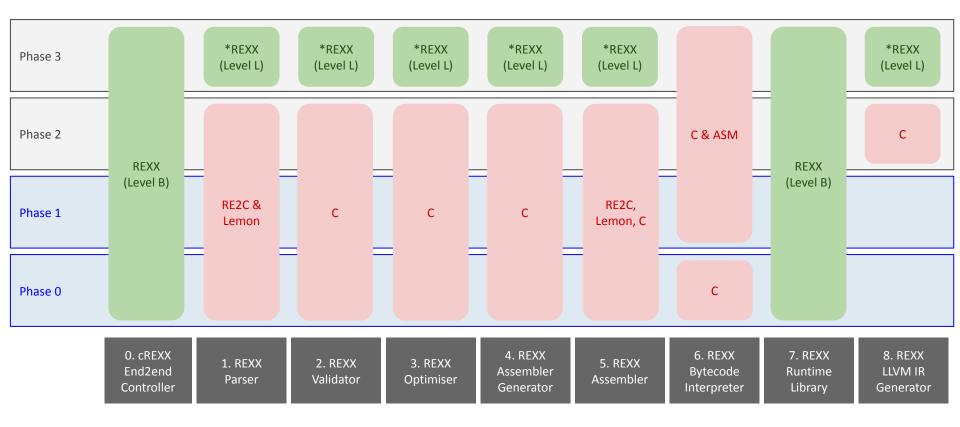






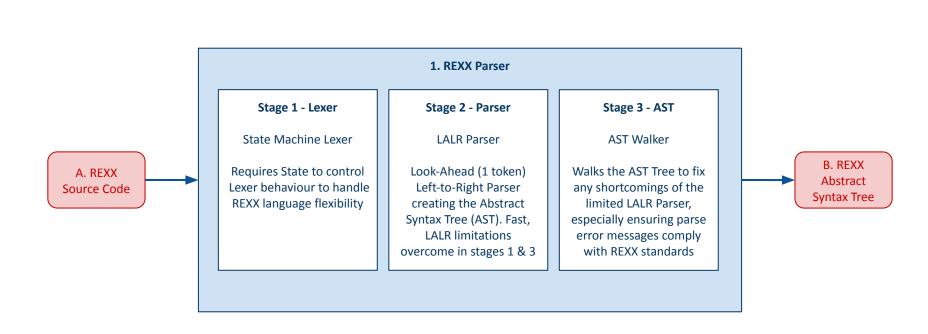


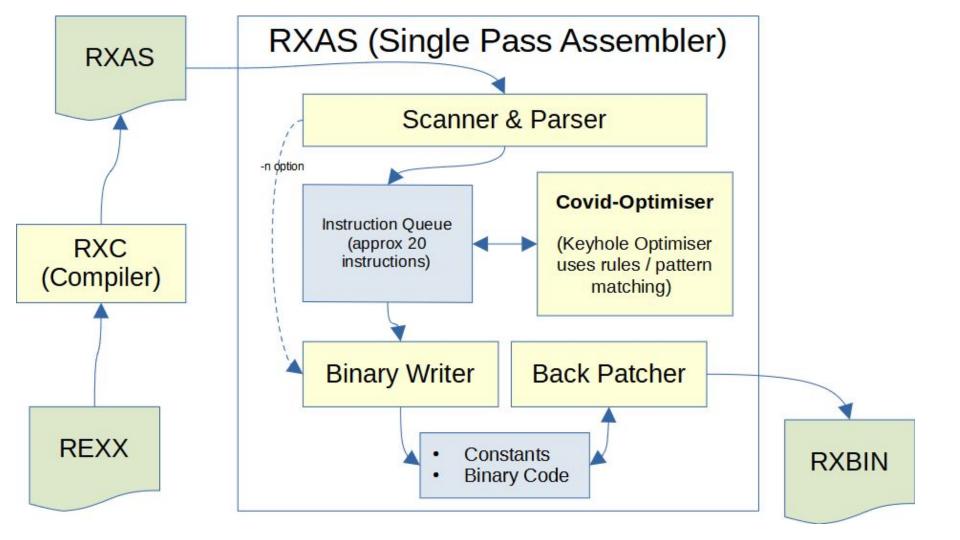
0. cREXX End2end Controller



- * REXX Level L provides the required:
 - 1. Extended PARSE to handle PEG Grammars
 - 2. Native support of Language Engineering data structures (ASTs and Symbol Tables)







cREXX in Action

How to Help?

How to Help?

- Github https://github.com/adesutherland/CREXX
- Contact myself or René
- Fortnightly Evening Zoom meetings

• Code - Test - Use - Feedback - or just lurk!

Thanks to ...

- René Jansen Our PM; for all his encouragement and work on the built in functions
- Peter Jacob Our microcode engineer!
- Mike Großmann For rolling up his sleeves when needed
- Michael Beer, Bob Bolch and everyone else who comes to our project meetings when they should be having a beer!

Adrian Sutherland

- CTO of Jumar Technology, specialists in legacy modernisation
- Journeyman Architect
- Keeps "hands-on" through numerous projects, from Raspberry PI toys and Domain Specific Languages to open architectural papers and other assets.

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